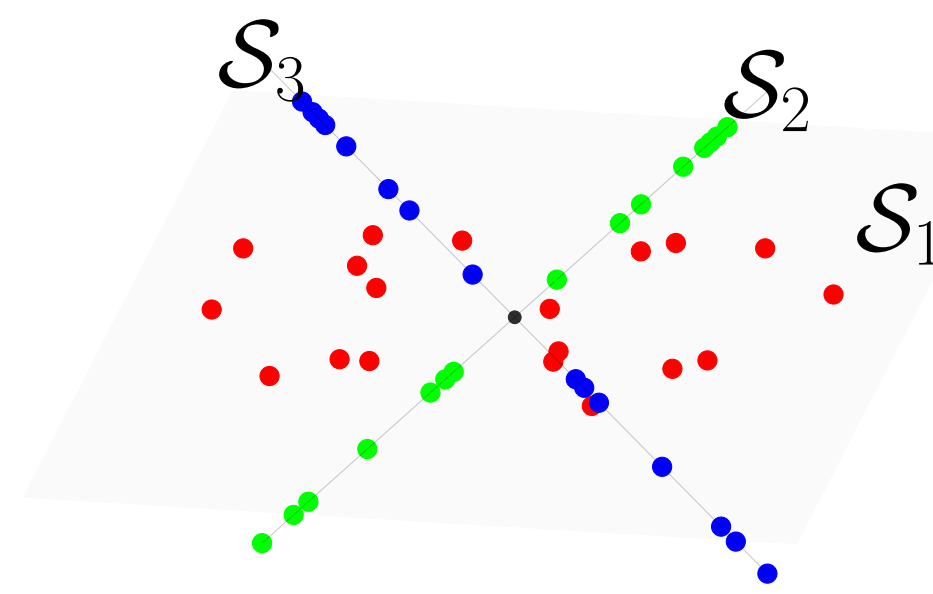


## Motivations

- **Sparse recovery:** Find the unique sparsest representation  $\mathbf{x}$  of a vector  $\mathbf{b}$  with respect to a dictionary  $\mathcal{A}$ :

$$\min_{\mathbf{x}} \|\mathbf{x}\|_0 \quad \text{s.t.} \quad \mathbf{b} = \sum_j \mathbf{a}_j x_j. \quad (1)$$

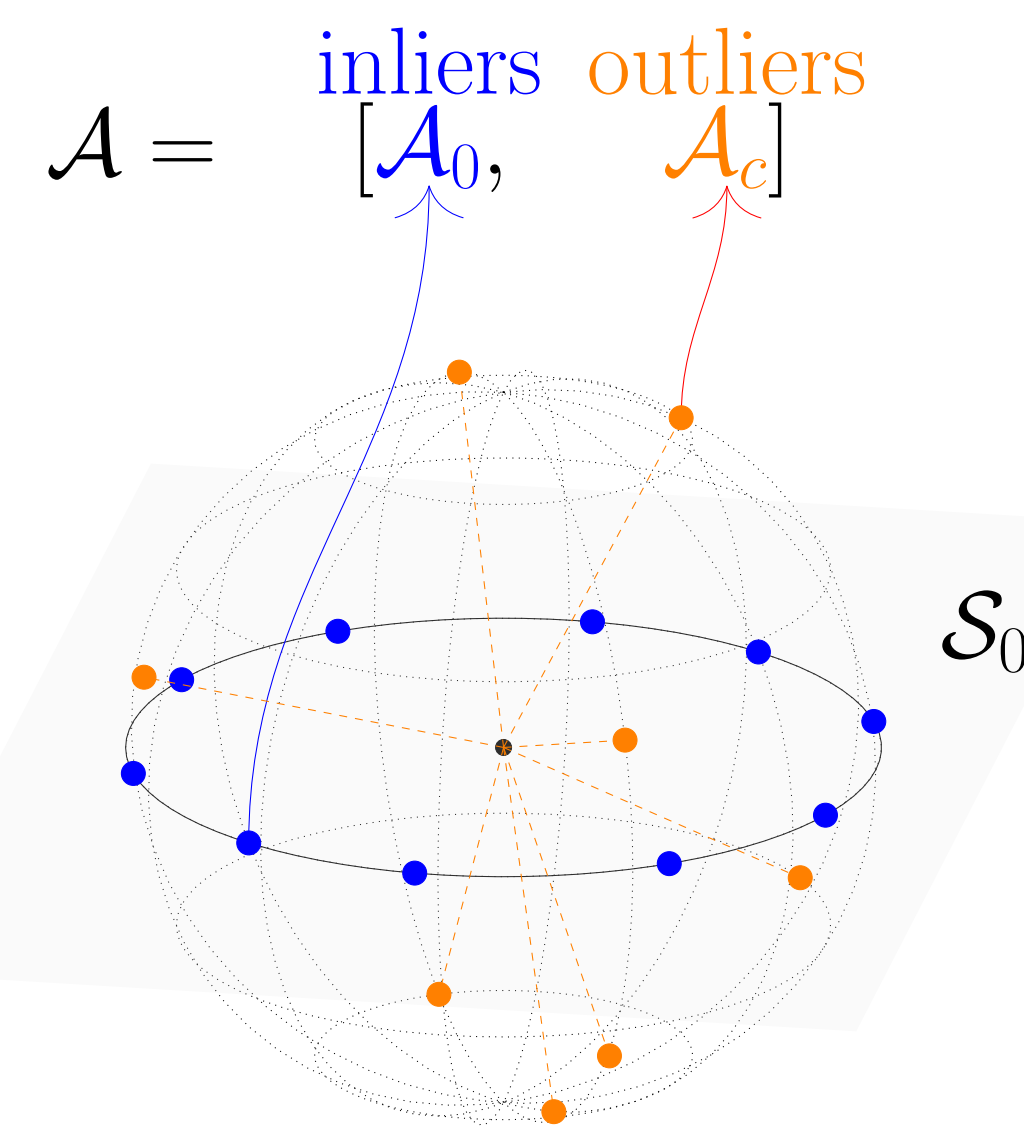


- **Subspace sparse recovery:** Find a subspace sparse representation  $\mathbf{x}$  of  $\mathbf{b}$  whose support gives its membership to one or more low-dimensional subspaces.
- **Applications:** in machine learning, computer vision and signal processing.

	Classification	Clustering
$\mathbf{b}$	Test data point	Data point
$\mathbf{a}_j$	All labeled data points	All other data points
$\mathbf{x}$	Gives membership	Gives pairwise affinity
Application	Face recognition	Face clustering

## Problem Statement

- **Subspace:**  $\mathcal{S}_0 \subset \mathbb{R}^D$ .
- **Subspace structured dictionary:**  
 $\mathcal{A} = \{\mathbf{a}_j \in \mathbb{R}^D, j \in J\} = \mathcal{A}_0 \cup \mathcal{A}_c$ ,  
 where  $\begin{cases} \mathcal{A}_0 \subseteq \mathcal{S}_0 : & \text{inliers to } \mathcal{S}_0 \\ \mathcal{A}_c \cap \mathcal{S}_0 = \emptyset : & \text{outliers to } \mathcal{S}_0 \end{cases}$



- **Subspace sparse representation:**

Let  $\mathbf{b} \in \mathcal{S}_0$ . Then  $\mathbf{x}_j \neq 0 \implies \mathbf{a}_j \in \mathcal{A}_0 \subseteq \mathcal{S}_0$ .

- **Problem:** Find conditions under which orthogonal matching pursuit (OMP) or basis pursuit (BP) give a subspace sparse representation for all  $\mathbf{b} \in \mathcal{S}_0$ .

## Challenges

- **Sparse recovery:** dictionary needs to have low mutual coherence or satisfy the restricted isometry property so that the sparsest solution be unique.
- **Subspace sparse recovery:**
  - Points in a subspace can be arbitrarily close, hence  $\mathcal{A}$  may not be incoherent.
  - Subspace sparse representations may not and need not be unique.

## Main Result: Geometric Conditions

- **Theorem:** The solutions of OMP and BP are subspace sparse for all  $\mathbf{b} \in \mathcal{S}_0$  if the dictionary  $\mathcal{A} = \mathcal{A}_0 \cup \mathcal{A}_c$  satisfies either of the following conditions:

$$\begin{aligned} \text{Principal Recovery Condition (PRC): } & \gamma_0 < s(\mathcal{A}_c, \mathcal{S}_0), \\ \text{Dual Recovery Condition (DRC): } & \gamma_0 < s(\mathcal{A}_c, \mathcal{D}_0). \end{aligned}$$

- **Notation:**

–  $s(\mathcal{A}_c, \mathcal{S}_0)$ : minimum angle between the outliers  $\mathcal{A}_c$  and the subspace  $\mathcal{S}_0$ , i.e.,

$$s(\mathbf{v}, \mathbf{w}) := \cos^{-1} \left\langle \frac{\mathbf{v}}{\|\mathbf{v}\|_2}, \frac{\mathbf{w}}{\|\mathbf{w}\|_2} \right\rangle, \quad s(\mathcal{V}, \mathcal{W}) := \inf_{\mathbf{v} \in \mathcal{V} \setminus \{0\}} \inf_{\mathbf{w} \in \mathcal{W} \setminus \{0\}} s(\mathbf{v}, \mathbf{w}).$$

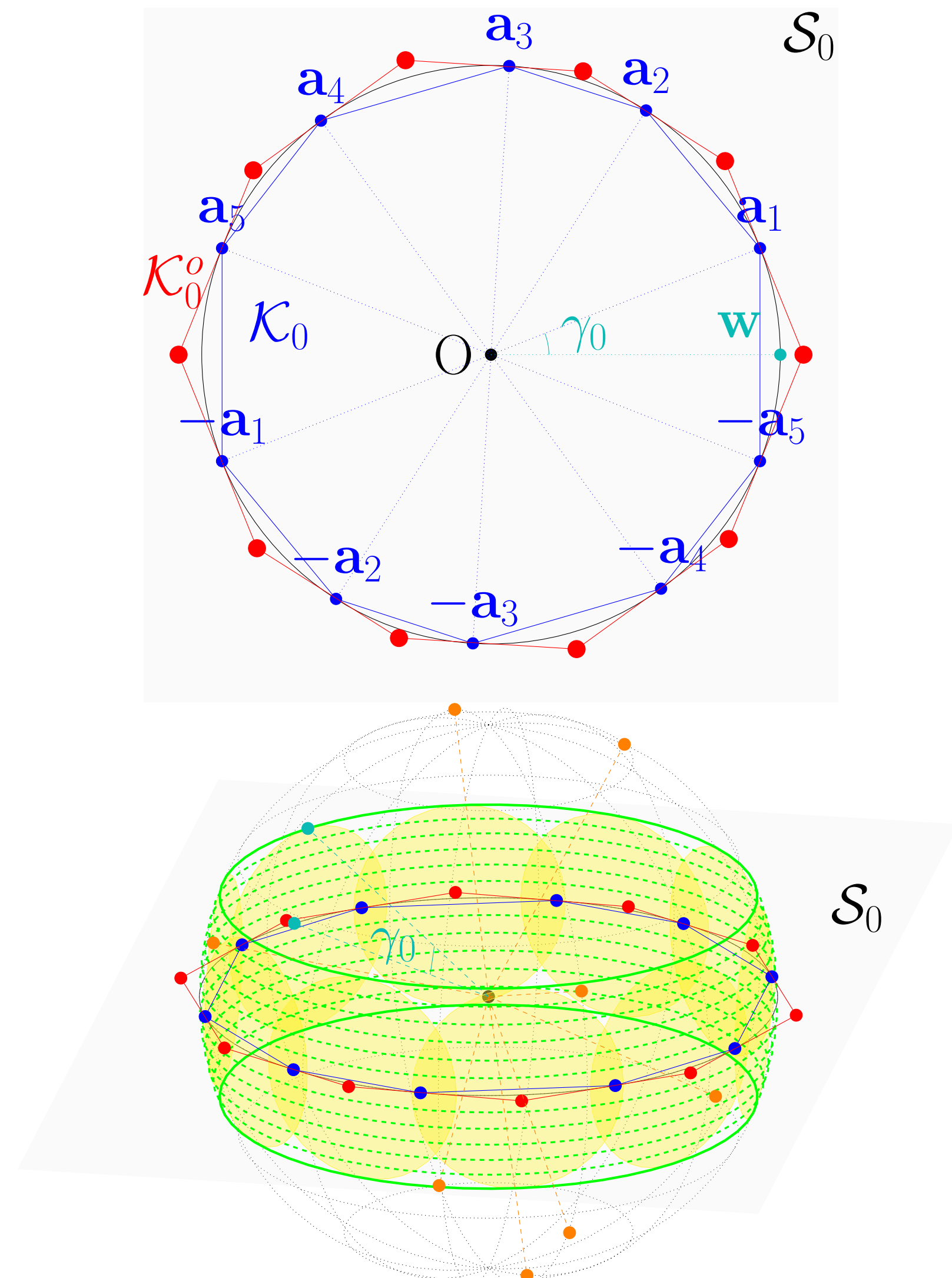
Note that  $s$  is larger if points from two sets are not close.

–  $s(\mathcal{A}_c, \mathcal{D}_0)$ : minimum angle between  $\mathcal{A}_c$  and  $\mathcal{D}_0$ , where  $\mathcal{D}_0 = \{\mathbf{v} \in \mathcal{S}_0 : \langle \mathbf{v}, \mathbf{w} \rangle \leq 1, \forall \mathbf{w} \in \text{conv}\{\pm \mathcal{A}_0\}\}$  is a finite subset of  $\mathcal{S}_0$ .

–  $\gamma_0 = \max\{s(\pm \mathcal{A}_0, \mathbf{w}) : \mathbf{w} \in \mathcal{S}_0\}$ : covering radius of points  $\pm \mathcal{A}_0$  in  $\mathcal{S}_0$ . Note that  $\gamma_0$  is smaller if the unit sphere of the subspace  $\mathcal{S}_0$  is densely populated with points  $\pm \mathcal{A}_0$ .

- **Geometric interpretation:**

– **PRC** holds if and only if all points in  $\mathcal{A}_c$  lie outside the **green region**.  
 – **DRC** holds if and only if all points in  $\mathcal{A}_c$  lie outside the **yellow region**.



## Implications for Sparse and Subspace Sparse Recovery

### Implications for subspace sparse recovery

- To achieve subspace sparse recovery, the inliers  $\mathcal{A}_0$  should be dense and well distributed in subspace  $\mathcal{S}_0$  so that  $\gamma_0$  is small, and the outliers  $\mathcal{A}_c$  should be sufficiently far from the subspace  $\mathcal{S}_0$  (by PRC) or the set  $\mathcal{D}_0$  (by DRC).
- DRC is a weaker condition than PRC since it only requires a finite subset of  $\mathcal{S}_0$  to be far from the outliers.
- PRC and DRC justify why subspace sparsity works better for low-dimensional subspaces, since the area of the PRC/DRC region relative to the area of the entire unit sphere is smaller.

### Implications for sparse recovery

- **Theorem:** Both OMP and BP recover any  $s$ -sparse signal with dictionary  $\mathcal{A}$  if  $\forall \mathcal{A}_0 \subseteq \mathcal{A}$  s.t.  $\text{card}(\mathcal{A}_0) = s$  we have  $\text{rank}(\mathcal{A}_0) = s$  and either PRC or DRC hold.
- **Properties:**
  - Geometrically interpretable: requires points to be “incoherent”.
  - Can be verified: explicit formula for computing  $\mathcal{D}_0$  and  $\gamma_0$ .
  - Less restrictive: the condition in theorem is implied by the mutual coherent condition  $\mu(\mathcal{A}) < \frac{1}{2s-1}$ , where  $\mu(\mathcal{A}) = \max_{j \neq k} |\langle \mathbf{a}_j, \mathbf{a}_k \rangle|$ .

## References

- [1] C. You and R. Vidal., Subspace Sparse Representation, In *Arxiv*, 2015.
- [2] M. Soltanolkotabi and E.J. Candes., A geometric analysis of subspace clustering with outlier, In *Annals of Statistics*, 2013.